

The Network Layer: Routing in the Internet

Dr. Michele Weigle

Department of Computer Science

Old Dominion University

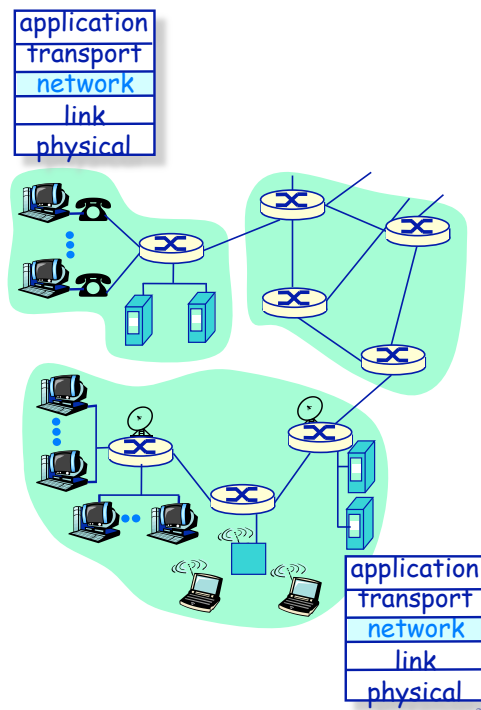
mweigle@cs.odu.edu

<http://www.cs.odu.edu/~mweigle/CS455-F08>

1

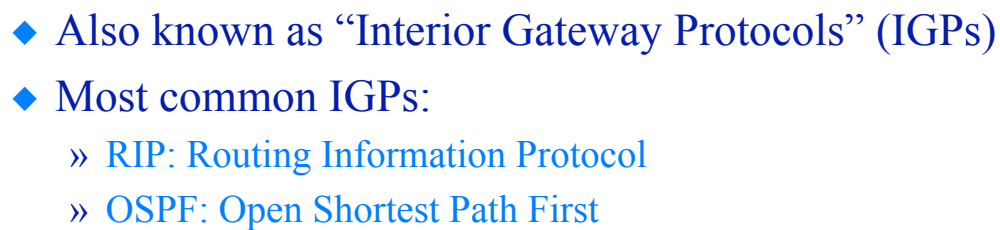
The Network Layer: Routing & Addressing Outline

- ◆ Network layer functions
- ◆ Virtual circuits and datagram networks
- ◆ Router architecture
- ◆ IP Internet Protocol
 - » Addressing
- ◆ Routing algorithms
 - » Least cost path computation algorithms
- ◆ Hierarchical routing
 - » Connecting networks of networks
- ◆ Routing on the Internet
 - » Intra-domain routing
 - » Inter-domain routing



2

Intra-AS routing

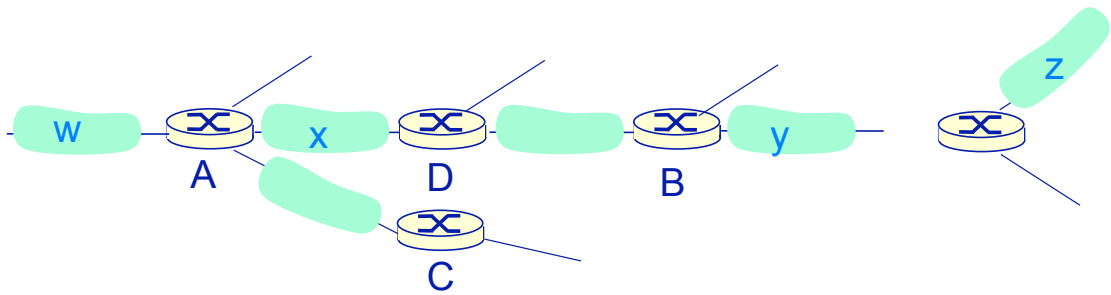


Routing Information Protocol (RIP)

- 4

RIP

Example



Routing Table in D

Destination Network	Next Router	Num. of hops to dest.
w	A	2
y	B	2
z	B	7
x	--	1
....

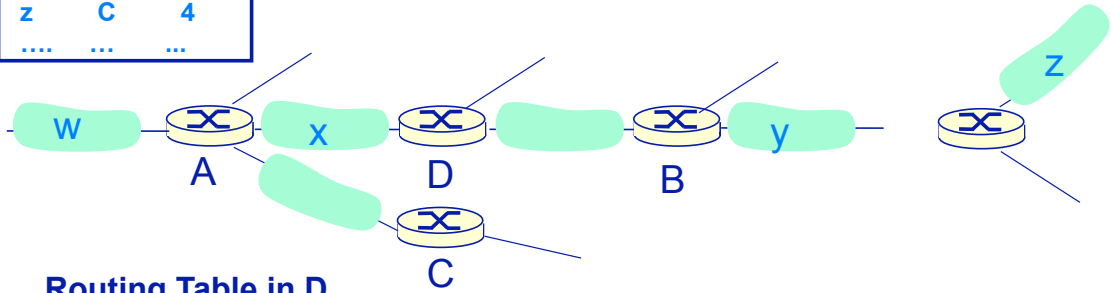
5

RIP

Example

Dest	Next	Hops
w	-	1
x	-	1
z	C	4
...

Advertisement from A to D



Routing Table in D

Destination Network	Next Router	Num. of hops to dest.
w	A	2
y	B	2
z	B A	7 5
x	--	1
....

6

RIP

Link failure and recovery

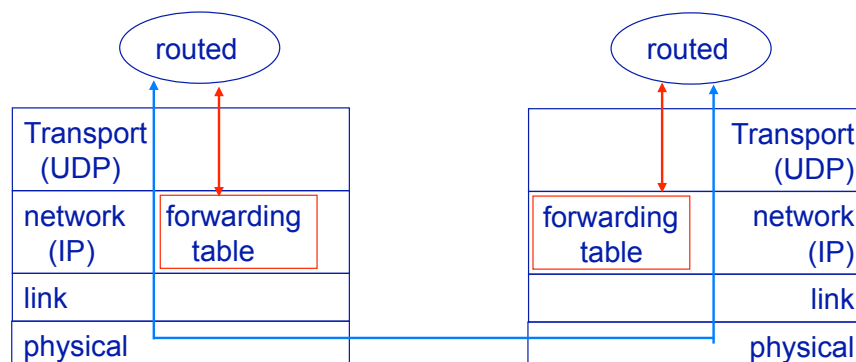
- ◆ If no advertisement heard after 180 seconds, adjacent node/link declared “failed”
 - » Routes via that adjacent node invalidated
 - » New advertisements sent to other adjacent nodes
 - » Advertisement receivers in turn send out new advertisements (if their tables changed)
 - » Link failure information quickly propagates to entire net
 - » Poisoned reverse used to prevent ping-pong loops
 - ❖ “infinity” = 16

7

RIP

Table processing

- ◆ RIP routing tables managed by application-level process called route-d (daemon)
- ◆ Advertisements sent in UDP packets (port 520), periodically repeated



8

Intra-AS routing

Open Shortest Path First (OSPF)

- ◆ “Open”: publicly available
- ◆ Uses the Link State minimum cost path computation algorithm
 - » LS update flooding
 - » Topology map at each node
 - » Route computation using Dijkstra’s algorithm
- ◆ OSPF advertisement carries one entry per adjacent node
- ◆ Advertisements disseminated to an entire AS (via flooding)

9

OSPF

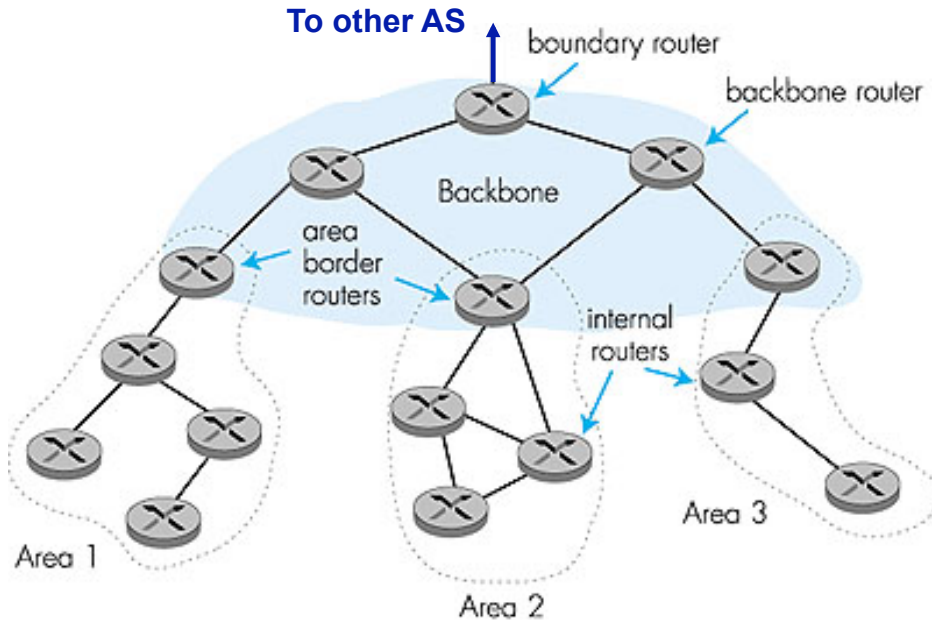
“Advanced” features (not in RIP)

- ◆ Security: all OSPF messages can be authenticated to prevent malicious intrusion
 - » TCP connections used in flooding
- ◆ Multiple same-cost paths can be used (only one path in RIP)
 - » packets in a connection may travel different paths even if routing tables don’t change
- ◆ For each link, multiple cost metrics for different network-layer “services”
 - » (e.g., satellite link cost set “low” for best effort; high for real time)
- ◆ Hierarchical OSPF used in large networks

10

OSPF

Hierarchical OSPF



12

OSPF

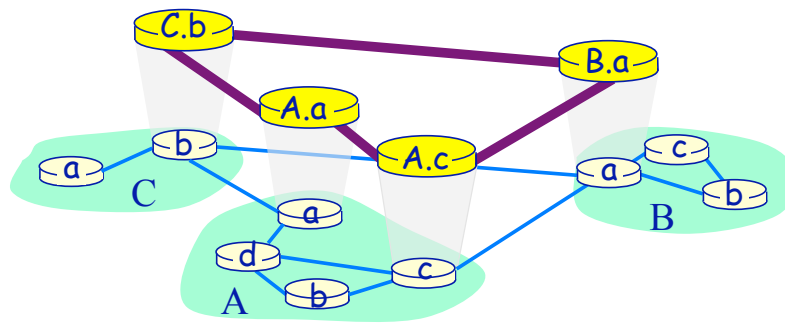
Hierarchical OSPF

- ◆ Two-level hierarchy: local area, backbone
 - » Link-state advertisements only in area
 - » Each area node has detailed area topology; only knows shortest path to networks in other areas
- ◆ Area Border Routers: “summarize” distances to nets in own area and advertise to other Area Border routers
- ◆ Backbone Routers: run OSPF routing limited to backbone
- ◆ Boundary Routers: connect to other AS

13

The Internet AS Hierarchy

Inter-AS Routing



- ◆ Border Gateway Protocol (BGP) is the *de facto* standard
- ◆ BGP allows each subnet to advertise its existence to the rest of the Internet (“I’m here!”)

14

BGP

Basics

- ◆ Pairs of routers (*BGP peers*) exchange routing information over TCP connections on port 179
 - » semi-permanent connection
 - » one connection per link that connects two gateway routers in different ASs (external BGP - eBGP)
 - » one connection per link that connects two gateway routers in same AS (internal BGP - iBGP)
- ◆ Allows each AS to learn which destinations are reachable via neighbors
 - » destinations are CIDR network prefixes

15

BGP

Path Attributes and BGP Routes

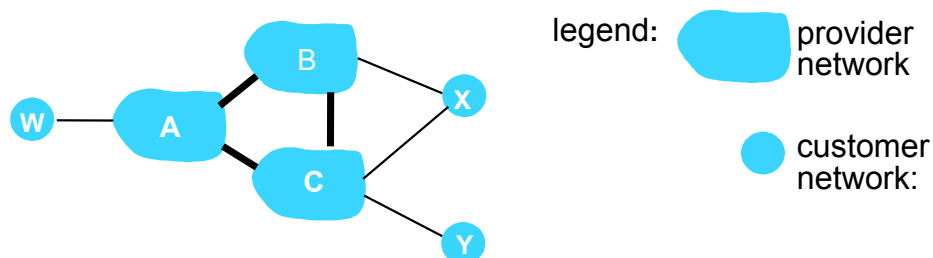
- ◆ Peers advertise entire routes
 - » AS-PATH
 - ❖ ASs through which the advertisement for the prefix has passed
 - ❖ when receiving an advertisement, each AS adds its ASN to the path list before sending its own advertisement
 - » NEXT-HOP
 - ❖ router interface that begins the AS-PATH

- ◆ Suppose gateway X sends a path to peer gateway W
 - » W may or may not select the path advertised by X
 - ❖ Cost, policy (“don’t route via competitor X’s network”), or loop prevention reasons
 - » If W selects the path advertised by X to Z, then:
$$path(W,Z) = W + path(X,Z)$$

16

BGP

Routing Policy

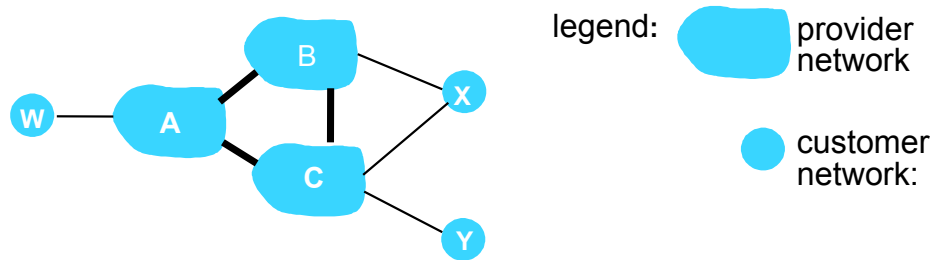


- ◆ A,B,C are *provider networks*
- ◆ X,W,Y are customer (of provider networks)
- ◆ X is *dual-homed*: attached to two networks
 - » X does not want to route from B via X to C
 - » .. so X will not advertise to B a route to C

17

BGP

Routing Policy



- ◆ A advertises path AW to B
- ◆ B advertises path BAW to X
- ◆ Should B advertise path BAW to C?
 - » No way! B gets no “revenue” for routing CBAW since neither W nor C are B’s customers
 - » B wants to force C to route to w via A
 - » B wants to route *only* to/from its customers!

18

The Internet AS Hierarchy

Why different intra- and inter-AS routing?

- ◆ Policy:
 - » Inter-AS: administration wants control over how its traffic routed and who routes through its network
 - » Intra-AS: single administration, so no “policy” decisions needed
- ◆ Scale:
 - » Hierarchical routing saves table size, reduced update traffic
- ◆ Performance:
 - » Intra-AS: can focus on performance
 - » Inter-AS: policy may dominate over performance

19

The Network Layer: Routing & Addressing

Outline

- ◆ Network layer functions
- ◆ Virtual circuits and datagram networks
- ◆ Router architecture
- ◆ IP Internet Protocol
 - » Addressing
- ◆ Routing algorithms
 - » Least cost path computation algorithms
- ◆ Hierarchical routing
 - » Connecting networks of networks
- ◆ Routing on the Internet
 - » Intra-domain routing
 - » Inter-domain routing

