Basic Components — Lexical

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3.1 Phases of Translation

Translators are divided into

- compilers, that produce machine code as output
- interpreters, that directly execute the source program



Lexical Analysis

source code

tokens

program output -

program input

In practice, almost all intepreters share at least some front-end phases with compilers.

Syntax

Analysis

Execution

Front End



3.2 Lexical

Characters occur in groups that have an "atomic" meaning in a language.

Such groups are called **tokens**. A language may have many different kinds of tokens.

The string of characters that corresponds to a given token is called a **lexeme**.

For example, in the code

if X > 1.5 THen

we have if tokens: identifier greater number then lexemes: if Х 1.5 THen > structure Semantic Tokens Analysis The kinds of tokens vary from one language to another, but some common ones are meaning • constants, operators, identifiers, keywords - If a keyword is not allowed to be used as an identifier, it is called a reserved word. In addition, some strings of characters don't contribute to any tokens: • white space comments Lexemes The strings that can make up a given kind of token will also vary between languages. E.g., identifiers X, longName, long_Name, \$name, name, NAME Semantic Analysis reserved words if, IF, iF, end, fi, endif constants 'abc', "abc", 0.275, .275, 0.275E3, 0.275G3 We describe the lexemes for a token kind either via grammars or via regular expressions. **Regular Expressions** In their simplest form, a **regular expression** R must be one of • a single character • ST, the concatenation of two other regular expressions • S|T, the choice of two regular expressions

Regular Expressions (cont.)

- ÷
- S*, 0 or more repetitions of a regular expression S
 - (known as the **Kleene closure**)
- (S), a regular expression within parentheses

Example: integers in most languages look sort of like this:

(+|-)(0|1|2|3|4|5|6|7|8|9)*

But this isn't quite right. Why not?

Common Extensions To Regular Expressions

- R+ denotes 1 or more repetitions of R
- R? denotes 0 or 1 occurrence of R
- R_m^k denotes between m and k occurrences of R

Common Extensions (cont.)

- [*abc*...] is short for (*a*|*b*|*c*|...), where only single characters can appear between the [].
 - The notation a z is also allowed within [], to denote a range of consecutive characters.
- Quote characters like '+' that would otherwise be confused.

With these, we can reduce our description of integer lexemes to

('+'|-)?[0-9]+

What do the following regular expressions describe?

- 0 * 1 *
- (0 * 1*)*
- (0 * 1 *) +
- (00|01|10|11)*

What do the following regular expressions describe?

- 0*1*
- (0 * 1*)*
- [01]+
- (00|01|10|11)*

What do the following regular expressions describe?

- 0 * 1*
- (0 * 1*)*
- [01]+
- (00|01|10|11)*

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- (0 * 1*)*
- [01]+
- (00|01|10|11)*

Here are some descriptions of identifier lexemes in different languages.

• Pascal:

[a - zA - Z][a - zA - Z0 - 9]*

- C: $[a zA Z_{-}][a zA Z0 9_{-}]*$
- FORTRAN: $[a zA Z][a zA Z0 9]_0^5$

Can you explain the differences?

What would be the lexemes for the reserved word "for" in C? in Pascal?

Not all lexical conventions can be described via regular expressions.

- For example, older languages such as FORTRAN and COBOL had column dependencies.
- ALGOL (published form) used typesetting information: if was an identifier, but **if** is a reserved word.